A Living System Must Have Noncomputable Models

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Abstract Chu and Ho’s recent article in *Artificial Life* is riddled with errors. In particular, they use a wrong definition of Robert Rosen’s *mechanism*. This renders their “critical assessment” of Rosen’s central proof null and void.

Keywords Robert Rosen, *Life Itself*, (M,R)-systems, noncomputability

1 Introduction

The conclusion of Chu and Ho’s recent article [1] published in *Artificial Life* is that “Rosen’s central proof is wrong.” The “central proof” refers to the main conclusion of Robert Rosen’s book *Life Itself* [5], that a living system is not a mechanism and consequently must have noncomputable models. Chu and Ho, however, use a definition of *mechanism* that is different from Rosen’s. This and numerous other errors in their article make their argument irrelevant. In this short note I shall discuss some of their errors.

To foreshadow what is to come, let me first point out that the category-theoretic definitions of product and coproduct (captions of Figures 1 and 3 in [1]) are wrong (and cannot be explained away as simple misprints). The proper definitions may be found in any introductory textbook in category theory. The definitive reference on this branch of abstract algebra is [4]. Products and coproducts are examples of universal constructions in category theory. Each universal determines a representation of a corresponding set-valued functor as a hom-functor. Rosen often used category theory as a metalanguage in his discussions, usually in the prototypical category Set, in which the objects are sets (without any further requisite structure) and the morphisms are mappings.

2 Rosen’s Theorems

In [5], Rosen defined the term *simulable* and several of its synonyms. A mapping is simulable if it is “definable by an algorithm.” It is variously called computable, effective, and evaluable by a mathematical (Turing) machine. In Chap. 8 of [5] he gave the following:

**Definition 2.1:** A natural system $N$ is a mechanism if and only if all of its models are simulable.

He then proved five propositions for a mechanism $N$. In particular, “Conclusion 4” is the following:

**Theorem 2.2:** Analytic and synthetic models coincide in the category $\mathbf{C}(N)$ of all models of $N$; direct sum $= \text{direct product}$.

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And “Conclusion 5” is the following:

**Theorem 2.3:** Every property of $N$ is fractional.

Immediately following this, in Chap. 9 of [5], Rosen, using these five just-proven properties, presented a detailed *reductio ad absurdum* argument that proves that certain modes of entailment are not available in a mechanism:

**Theorem 2.4:** There can be no closed path of efficient causation in a mechanism.

The contrapositive statement of Theorem 2.4 is

**Theorem 2.5:** If a closed path of efficient causation exists in a natural system $N$, then $N$ cannot be a mechanism.

Taking Definition 2.1 of mechanism into account, this is equivalent to

**Theorem 2.6:** If a closed path of efficient causation exists for a natural system $N$, then it has a model that is not simulable.

An iteration of “efficient cause of efficient cause” is inherently hierarchical. A closed path of efficient causation must form a hierarchical cycle. Both the hierarchy and the cycle (closed loop) are essential attributes of this closure.

In formal systems, hierarchical cycles are manifested by impredicativities, or the inability to replace these self-referential loops with finite syntactic algorithms. Impredicativities are simply part of the semantic legacy of mathematics as a language, in their expression of transcendental operations. Further elaboration on this concept may be found in abundance in [6]. The nonsimulable model in Theorem 2.6 contains a hierarchical closed loop that corresponds to the closed path of efficient causation in the natural system being modeled. In other words, it is a formal system with an impredicative loop of inferential entailment. Thus we also have:

**Theorem 2.7:** If an impredicative loop of inferential entailment exists for a formal system, then it is not simulable.

A natural system that has a nonsimulable model is defined by Rosen as a complex system (Chap. 19 of [6]). A necessary condition for a natural system to be an organism is that it is closed to efficient causation (Chap. 1 of [6]). Theorem 2.7 then says an organism must be complex. The implication on the concept of artificial life is this:

**Theorem 2.8:** A living system must have noncomputable models.

All Rosen’s theorems have been mathematically proven (although Rosen’s presentations are not in the ordinary form of definition-lemma-theorem-proof-corollary that one finds in conventional mathematics journals). Indeed, no logical fallacy in Rosen’s arguments has ever been demonstrated. Counterexamples cannot exist for proven theorems. For a detailed exposition of the underlying logic, the reader is encouraged to consult [2].

Note that Rosen’s conclusion is not that artificial life is impossible. It is, rather, that life is not computable: However one models life, natural or artificial, one cannot succeed by computation alone. Life is not definable by an algorithm. There is, indeed, practical verification from computer science that attempts at implementation of a hierarchical closed loop lead to deadlock, and hence are forbidden in systems programming [7].
3 Set Theory

Let $\mathcal{N}$ be the collection of all natural systems. By Rosen’s Definition 2.1, the set of all mechanisms is

$$
\mathcal{M} = \{N \in \mathcal{N} : \text{all models of } N \text{ are simulable}\}.
$$

Let me also define

$$
\mathcal{Q} = \{N \in \mathcal{N} : \text{in the category } \mathcal{C}(N) \text{ of models of } N \text{ the collections of analytic models and synthetic models coincide}\}.
$$

and

$$
\mathcal{R} = \{N \in \mathcal{N} : \text{every property of } N \text{ is fractionable}\}.
$$

Then Theorem 2.2 says that $\mathcal{M} \subseteq \mathcal{Q}$, while Theorem 2.3 says that $\mathcal{M} \subseteq \mathcal{R}$.

It is important to note that Rosen's five properties for a mechanism $N$ are necessary properties. If $N$ is a mechanism, then $N$ necessarily has each one of these properties. Rosen only needed the necessity in these statements to establish the subsequent theorems. The five conclusions do not say the converse, that if $N$ has any one of these properties, then it is sufficient to guarantee that $N$ is a mechanism. In particular, note that while $\mathcal{M} \subseteq \mathcal{Q}$, we in fact have $\mathcal{M} \neq \mathcal{Q}$.

Rosen himself fully realizes that $\mathcal{M}$ is a proper subset of $\mathcal{Q}$. On p. 186 of [5] he wrote “EVERY MODEL SIMULABLE IMPLIES ANALYTIC=SYNTHETIC and, to a sufficiently large extent, conversely.” This says $\mathcal{M}$ is “almost” all of $\mathcal{Q}$, but it is possible to find counterexamples of natural systems in $\mathcal{Q}$ but not in $\mathcal{M}$.

The fatal error in the Chu and Ho’s article [1] is that they misrepresent Rosen’s definition of mechanism:

Rosen defines mechanisms as the class of systems of which all analytic models and synthetic models are equivalent (first sentence of Section 2.4 of [1])

and then proceed to define mechanism erroneously:

“we will give an alternative, yet essentially equivalent definition of mechanism: A system is a mechanism if all its analytic models are equivalent to synthetic models” (second paragraph in Section 3 of [1])

In other words, the set of mechanisms in Chu and Ho’s definition is $\mathcal{Q}$, not $\mathcal{M}$.

The main argument in Chu and Ho’s supposed demonstration of “why Rosen’s central proof is wrong” appears to be their construction of a system with equivalent analytic and synthetic models (which they erroneously identified as a mechanism) that is nevertheless not fractionable. Their contention is that this provides a counterexample to Theorem 2.3, and therefore by gross generalization all of Rosen’s results are suspect. They may or may not have an example of an element in $\mathcal{Q}$ but not in $\mathcal{R}$, but that is irrelevant: Theorem 2.3 says $\mathcal{M} \subseteq \mathcal{R}$, not $\mathcal{Q} \subseteq \mathcal{R}$. From wrong definitions arise nonsensical conclusions.
Let me elaborate on what I mean by “wrong definition.” Chu and Ho may of course define “mechanism” any way they please. After all, as Humpty Dumpty said to Alice: “When I use a word, it means just what I choose it to mean—neither more nor less.” But from a non-Rosen definition of mechanism one cannot expect a Rosen property of mechanism to automatically follow. I again urge the reader to consult [2] to understand why the search for counterexamples to Rosen’s theorems is futile.

4 (M,R)-Systems

Rosen devised a class of relational models of organisms called (M,R)-systems. He has discussed them on numerous occasions, including Section 10C of [5] and Chap. 17 of [6]. I have written on their noncomputability and realizations in two recent articles [2, 3]. The reader may refer to any or all of the above for their details, which need not be repeated here.

For our present purpose, we only need to know that an (M,R)-system is an example of a formal system that has an impredicative loop of inferential entailment, and hence is nonsimulable. Chu and Ho attempt a discussion of (M,R)-systems in Section 4 of [1], but get it wrong as well.

It is telling that they claim “Wolkenhauer [8] offers a very clear discussion” on (M,R)-systems, yet they still manage to misinterpret Wolkenhauer. So it would appear, unfortunately, that Chu and Ho do not correctly understand the concept of the replication map in (M,R)-systems, as explained by either Rosen or Wolkenhauer. I shall use Chu and Ho’s symbolism in my explanation in the following, although it is nonstandard and awkward.

There are three mappings in an (M,R)-system on three hierarchical levels, and they entail one another in a cyclic permutation. They are

- metabolism $\mathbf{f} : A \rightarrow B$
- repair $F : B \rightarrow \mathbf{f}$
- replication $B : \mathbf{f} \rightarrow F$

The cyclic entailment pattern when one combines these three maps, $\{\mathbf{f}, F, B\}$, is the closed hierarchical loop of inferential entailment in the (M,R)-system. Note that replication is $B : \mathbf{f} \rightarrow F$; not $\mathbf{f} : B \rightarrow F$ as claimed by Chu and Ho. The correct relational diagram in graph-theoretic form is not Figure 7 of [1], but Figure 1 of this letter (again using Chu and Ho’s symbolism).

5 Artificial Life

Let me reiterate the fact that Rosen did not say that artificial life is impossible, only that life is noncomputable. Indeed, the subtitle of [5] is “A comprehensive enquiry into the nature, origin, and..."
fabrication of life,” with a positive elaboration on the last topic. Artificial life does not have to be
limited to what a computing machine can do algorithmically. The first step is to admit that not
everything is computable, that is, throw away the Cartesian and Newtonian machine metaphor.
One must loosen the mechanist constraints and assert the existence of natural systems with
nonsimulable models.

Artificial life is not simply a simulation of life; it needs to be a model of life. The difference between
a simulation and a model is that in the latter the morphism is mapped along with the domain and
codomain, making the modeling relation a natural transformation in category-theoretic terms. (See
Chap. 7 of [5] for details.) Artificial life must have entailment patterns that are congruent with the
entailment patterns of living systems. These entailment patterns contain impredicative loops within
themselves, and are beyond finite syntactic computation. A thorough discussion on the fabrication of
life is found in Chap. 17 of [6], entitled “What does it take to make an organism?” Let me close by
quoting a paragraph from it.

On these grounds, we can see that the fabrication of something (e.g., an organism) is
a vastly different thing than the simulation of its behaviors. The pursuit of the latter
represents the ancient tradition that used to be called biomimesis, the imitation of life.
The idea was that by serially endowing a machine with more and more of the simulacra
of life, we would cross a threshold beyond which the machine would become an organism.
The same reasoning is embodied in the artificial intelligence of today, and it is articulated
in Turing’s Test. This activity is a sophisticated kind of curve-fitting, akin to the assertion
that since a given curve can be approximated by a polynomial, it must be a polynomial.

References
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