Primer design using genetic algorithm
Jain-Shing Wu¹,*, Chungnan Lee¹, Chien-Chang Wu³ and Yow-Ling Shiu²

¹Department of Computer Science and Engineering and ²Institute of Biomedical Sciences, National Sun Yat-Sen University, Kaohsiung, Taiwan and ³Department of Food Engineering, Da-Yeh University, Taiwan

Received on September 18, 2003; revised on December 12, 2003; accepted on January 19, 2004
Advance Access publication February 26, 2004

ABSTRACT
Motivation: Before performing a polymerase chain reaction experiment, a pair of primers to clip the target DNA sub-sequence is required. However, this is a tedious task as too many constraints need to be satisfied. Various kinds of approaches for designing a primer have been proposed in the last few decades, but most of them do not have restriction sites on the designed primers and do not satisfy the specificity constraint.
Results: The proposed algorithm imitates nature’s process of evolution and genetic operations on chromosomes in order to achieve optimal solutions, and is a best fit for DNA behavior. Experimental results indicate that the proposed algorithm can find a pair of primers that not only obeys the design properties but also has a specific restriction site and specificity. Gel electrophoresis verifies that the proposed method really can clip out the target sequence.
Availability: A public version of the software is available on request from the authors.
Contact: wujs@mail.cse.nsysu.edu.tw

INTRODUCTION
With the growing number of researches in bioinformatics, many biotechnologies have improved considerably. However, most of them need to amplify the number of DNA sequences, since the experimental results of the currently used biotechnologies are not very clearly recognized by the naked eye when only a small amount of DNA is used. Consequently, DNA cloning is becoming increasingly important. Polymerase chain reaction (PCR) is a method for the fast mass duplication of DNA sequences (Mullis and Faloona, 1987). Prior to performing a PCR, a pair of sub-sequences of DNA called ‘primers’ must be found in order to clip the target in a long DNA sequence.

Due to the properties of oligonucleotides that influence the efficiency of the PCR amplification, the optimal primer design includes criteria (McPherson et al., 1993; Sambrook and Russell, 2001) such as melting temperatures, length, base composition, 3’-termini, repeated and self-complementary sequences and complementarity between members of a primer pair. To ensure the primers will be efficiently annealed during each cycle of the PCR, the calculated values for the melting temperature of a primer pair should not differ by more than 5°C. The lengths of the members of a primer pair should not differ by more than 3 bp. The GC content of the members of a primer pair should be between 40 and 60%. The nature of the 3’ end of the primers is crucial. If possible, the 3’ end of each primer should be G or C. Actual differences for these criteria are aggregated by weighting sums. The prime design is to construct candidates and to select the best.

Various kinds of approaches in designing primers have been proposed in the last few decades. The manual primer design method can find a primer that fits the primer design constraints. However, it is too time consuming to find a good primer. Besides, accuracy can easily be lost through human errors. Since experiments are expensive, and a minor mistake may cause the experiment to fail, the manual primer design method is considered to be potentially unstable.

To improve the efficiency in primer design, many primer design software have been developed for DNA sequences already identified. However, if the designed primer does not satisfy a user’s requirements, then it is up to the user to modify the primers in order to meet the requirements. The ‘GeneFisher’ system proposed by Meyer et al. (1995, http://bibiserv.techfak.uni-bielefeld.de/docs/gf_paper.html) first aligns the unknown DNA sequence with the known DNA sequence, which has the same functions, and then designs the primer that can deal with the unknown DNA sequence. The system Consensus-DEgenerate Hybrid Oligonucleotide Primer (CODEHOP) proposed by Rose et al. (1998) can find primers for known amino acid sequences. It first calls upon a sequence alignment program, such as FASTA, to find a similar sequence. Then it retransforms the amino acid sequences into DNA sequences, and designs the primer for the similar DNA sequences. Although these systems have provided a convenient way for primer design, most of them do not have restriction sites on the designed sequences.
primers. Singh et al. (1998) proposed the system ‘Primer Premier’, which provides a good performance on primer designs. It provides a restriction enzyme graph for the DNA sequences by finding the positions of the restriction sites on the aligned DNA sequences. According to the graph, it designs the primer that contains specific restriction sites within the primers. However, it sometimes cannot find the solution, due to the fact that design properties may not be followed. Recently, Kämpke et al. (2001) employed dynamic programming to design primers. Their algorithm can solve the situation of designing multiple primers for multiple target DNA sequences. But, they did not consider some necessary criteria of primer design, such as avoiding ‘T’ or ‘A’ at the 3’ end. Besides, in order to reduce computational complexity, the algorithm removes some potential solutions that may take a long time to compute.

Since the search space of the primer design problem is huge and complex, a better method to solve this problem is needed. Genetic algorithms (GAs) are well-known heuristic algorithms based on the imitation of natural systems. Their effectiveness in search and optimization problems has received extensive attention. GAs imitate nature’s evolutionary process and the genetic operations on chromosomes in order to achieve optimal solutions. In each run of a search, it generates a new and usually a better generation of solutions than the previous run (Goldberg, 1989; Jong, 1988). Due to the nature of the GA process being similar to the evolution of DNA, it is suitable for solving the primer design problem.

In this paper, a new algorithm using 9 GA for designing primers for PCR is presented. Since most of the design properties are treated as fitness rules, solutions that do not obey the fitness rule are eliminated through competition. The proposed algorithm also searches for restriction sites and specificity. If the restriction sites are closed to a user’s requirements, then the rank of this primer is promoted so that it can be listed on the top of the solution ranking set.

The rest of this paper is organized as follows. In the next section, we present some essential definitions for the proposed algorithm. In the following section, we describe the process of the proposed algorithm for PCR primer design. Experimental results are given in the succeeding section. Discussions are given in fifth section and finally conclusions are drawn in the last section.

DEFINITION OF THE PROPOSED ALGORITHM

Let $G_D$ be the DNA sequence template, which is denoted as the template of the base-nucleic acid code sequence of DNA. For example, $G_D$ is represented as

$$G_D = AATCGACCACCT...$$

where A, T, C and G are the base-nucleic acid codes. $\overline{G_D}$ is denoted as the complement code of the original base-nucleic acid code. The complement of A is T and vice versa. Similarly, C and G are the complements of each other. For example, $G_D$, which was described above, its complement $\overline{G_D}$ is

$$\overline{G_D} = TTAGCTGGTAT...$$

The forward primer of $G_D$ is denoted as $B_f$, and is defined as follows:

$$B_f = \{b_i \mid i \text{ is the index of } G_D \text{ between } F_s \text{ and } F_e\},$$

where $F_s$ and $F_e$ are, respectively, denoted as the start index and the end index of $B_f$ in $G_D$. The reverse primer of $G_D$ is denoted as $B_r$, and is defined as follows:

$$B_r = \{b_i \mid i \text{ is the index of } \overline{G_D} \text{ between } R_s \text{ and } R_e\},$$

where $R_s$ and $R_e$ are, respectively, denoted as the start index and the end index of $B_r$ in $\overline{G_D}$. The individual of the proposed algorithm is denoted as one pair of primers, which is presented as a vector $P_t$, and is written as

$$P_t = (F_s, F_e, R_s, R_e).$$

Since the four components of the individual are dependent, if one of them is changed by the crossover or mutation process, it may sometimes cause an error, violating the length constraint. To avoid this problem, we transform the dependence of these four components into an independent form. Hence, the individual is transformed from a dependent form $(F_s, F_e, R_s, R_e)$ into an independent form $(F_s, \alpha, \beta, \gamma)$, and is defined as

$$P_t' = (F_s, \alpha, \beta, \gamma),$$

where $\alpha, \beta$ and $\gamma$ are given as follows:

$$\alpha = (F_e - F_s), \quad \beta = (R_s - F_e), \quad \gamma = (R_e - R_s).$$

The dependent form represents the actual position of the primer, and the independent form represents the relative position. For example, the dependent form of individual $P_t$ is (145, 164, 989, 1011). Therefore, the independent form $P_t'$ is (145, 19, 825, 22). Since the dependence of the four components are removed, one of the four components being changed does not cause a violation. Suppose $P_t$ is a primer of $G_D$, then the length $|P_t|$ is the sum of the numbers of all base-nucleic acid codes, and is written as

$$|P_t| = \#G + \#C + \#A + \#T.$$
the Wallace formula, is written as

$$T_m(P_1) = (\#G + \#C) \times 4 + (\#A + \#T) \times 2.$$  

The GC content, $GC(P_1)$, is the ratio of base-nucleic acid codes G and C of sequence $P_1$. $GC(P_1)$ is written as

$$GC(P_1) = \frac{\#G + \#C}{|D_A|} \times 100\%.$$  

The specificity of primer $P_1$ is denoted as $Uni(P_1)$, which is to examine the annealing position of the primer $P_1$ in $GD$, and it is

$$Uni(P_1) = \begin{cases} 0, & \text{if } P_1 \text{ appear in } GD \text{ once,} \\ 1, & \text{if } P_1 \text{ appear in } GD \text{ more than once.} \end{cases}$$  

The termination of primer $P_1$ is denoted as $Term(P_1)$, which is to examine whether the 3′ end of the primer $P_1$ is G or C, and is defined as

$$Term(P_1) = \begin{cases} 0, & \text{if } 3' \text{ end is } G, C, CG \text{ or GC,} \\ 1, & \text{otherwise.} \end{cases}$$  

**THE PROPOSED ALGORITHM FOR THE PCR PRIMER DESIGN**

The proposed algorithm for the PCR primer design consists of the initialization process, evaluation process, crossover process and the mutation process. Figure 1 shows the flowchart of the proposed algorithm and detailed steps of the proposed GA for the PCR primer design.

The initialization process randomly generates the initial individuals of population $P$, which is a set of 500 solutions $P_t$. In this phase, the $F_s$ of one forward $B_f$ in $GD$ is generated randomly. Then, $\alpha$ and $\gamma$ are generated randomly within $18–26$; $\beta$ is thus generated randomly around the length of the target product. Then, variables $F_s$, $\alpha$, $\beta$ and $\gamma$ are recorded in the matrix, in order to avoid the duplicate primer pair from being generated in the population.

The evaluation process checks whether the solutions satisfy the design constraints or not. When solution $P_t$ is generated, the length of each primer $P_t$ should be within $18–26$. If the length is <18, it is not easy to control the melting temperature. Besides, if one wants to optimize the melting temperature, then the primer must become GC-rich, which may cause a violation of specificity. Suppose that $B_f$ is the forward primer of $P_t$, and $B_r$ is the reverse primer of $P_t$. The length check function is denoted as $leng(P_t)$ and is defined as

$$leng(P_t) = \begin{cases} 0, & \text{if } 18 \leq |B_f|, |B_r| \leq 26, \\ 1, & \text{otherwise.} \end{cases}$$  

Besides, the length different function lengd($P_t$) between $B_f$ and $B_r$ must be <3mer, and it is

$$lengd(P_t) = \begin{cases} 0, & \text{if } ABS(|B_f| - |B_r|) \leq 3, \\ 1, & \text{otherwise,} \end{cases}$$  

where $ABS(.)$ denotes the absolute value. The evaluation process evaluates population $P$ based on the design properties that are considered as constraints. The first constraint, the melting temperature, ensures that the primer pair will smoothly anneal to $GD$ at the same temperature. The temperature difference $Tmd$ between $B_f$ and $B_r$ is defined as

$$Tmd(P_t) = \begin{cases} 0, & \text{if } ABS[Tm(B_f) - Tm(B_r)] \leq 5, \\ 1, & \text{otherwise.} \end{cases}$$  

The GC content check function $GC_{P_t}$ of $P_t$ is denoted as

$$GC_{P_t} = \begin{cases} 0, & \text{if } 40\% \leq GC(B_f), GC(B_r) \leq 60\%, \\ 1, & \text{otherwise.} \end{cases}$$  

The specificity of the primers should be zero:

$$Uni(B_f) = 0, \quad Uni(B_r) = 0, \quad Uni(P_t) = Uni(B_f) + Uni(B_r).$$  

The termination of the primers should also be zero:

$$Term(B_f) = 0, \quad Term(B_r) = 0, \quad Term(P_t) = Term(B_f) + Term(B_r).$$  

The proposed algorithm for the PCR primer design consists of the initialization process, evaluation process, crossover process and the mutation process. Figure 1 shows the flowchart of the proposed algorithm and detailed steps of the proposed GA for the PCR primer design.
sequences of one primer should not be able to bind to any site on the other primer. Self-complementary sequences of primers should be avoided:

\[
\text{Sc}(P_t) = \begin{cases} 
0, & \text{if there is no self-complementary} \\
& \text{of } B_f \text{ and } B_r, \\
1, & \text{otherwise.}
\end{cases}
\]

In a primer pair, one primer should not be the complement of the other one:

\[
\text{PC}(P_t) = \begin{cases} 
0, & \text{if there is no pair-complementary} \\
& \text{of } B_f \text{ and } B_r, \\
1, & \text{otherwise.}
\end{cases}
\]

Since the constraint termination should be obeyed, but because three or more Cs or Gs at the 3’ ends of primers may promote error annealing at G- or C-rich sequences, this condition should be avoided. Restriction sites are situated at the primer if possible. The proposed algorithm searches the individual to see if there is a similar restriction site on it. The typical restriction sites used in this paper are listed in the following table:

<table>
<thead>
<tr>
<th>Restriction enzyme</th>
<th>Sequence</th>
</tr>
</thead>
<tbody>
<tr>
<td>ApaI</td>
<td>GGGCCC</td>
</tr>
<tr>
<td>AvrII</td>
<td>CCTAGG</td>
</tr>
<tr>
<td>BamHI</td>
<td>GGATCC</td>
</tr>
<tr>
<td>BglII</td>
<td>AGATCT</td>
</tr>
<tr>
<td>DraI</td>
<td>TTTAAA</td>
</tr>
</tbody>
</table>

In order to recruit a specific restriction site, the proposed algorithm first checks the individual from 5’ end to 3’ end for a proceed pattern match, and verifies whether the enzyme is on it or not. If there is a similar restriction site (the matched pattern’s length \(|P_m|\) is less or equal to the length of the enzyme \(L_e\) and more or equal to the length of the enzyme minus 3, \(|(L_e - 3) \leq |P_m| \leq L_e|\), then the proposed algorithm adjusts the fitness value to allow the individual most likely to pass the evaluation. The restriction site check function \(R_t(P_t)\) is denoted as

\[
R_t(P_t) = \begin{cases} 
0, & \text{if there exists a restriction site} \\
& \text{of } B_f \text{ or } B_r, \\
1, & \text{otherwise.}
\end{cases}
\]

In addition to the examination of the restriction site, the other time-consuming activity is specificity. For specificity there is an attempt to find an actual position for a primer pair to anneal to. If there are several positions in the target that are not mismatched to other positions of the target sequence for the primer pair to anneal to, then the product will consist of non-specific amplified sequences, so that the result affects the experiments that follow. Due to the time-consuming process of examining all individuals in order to reduce the possibility of mis-priming, we use a matrix, which records the position and the specificity fitness value of the position, to speed up the examination process. Based on the simulation, with the use of the matrix the algorithm can run four to five times faster than the algorithm without using the matrix.

The fitness value is computed by the design constraints mentioned above, and is as follows:

\[
\text{Fitness}(P_t) = \text{leng}(P_t) + 3 \ast \text{lengd}(P_t) + 3 \ast \text{Tmd}(P_t) + 3 \ast \text{GC}(P_t) + 3 \ast (\text{Term}(B_f) + \text{Term}(B_r)) + 50 \ast \text{Uni}(P_t) + 10 \ast \text{Sc}(P_t) + 10 \ast \text{PC}(P_t) + R_t(P_t)
\]

Selection applies the Roulette Wheel method to allow the individuals with a high weight to have a higher chance to be selected, and sends these two individuals into the mating pool. The weight mentioned here is the inverse of the fitness value, which is calculated in the evaluation process. Figure 2 shows the flowchart of the crossover process and the mutation process.

The crossover process generates a random number \(R\) that is smaller than 16. It uses the binary form of \(R\) as a mask to decide which components of individuals \(X\) and \(Y\) should be exchanged. For example, a random number \(R\) is 11(10) and its binary is 1011(2). The first, third and fourth components of individuals \(X\) and \(Y\) should be exchanged. Then, the crossover process examines the offspring individuals as to whether the offspring violate the constraints or not. Figure 3 shows an example of the crossover process.

The four components of \(X\) are represented as \((145, 19, 825, 22);\) and the four components of \(Y\) are represented as \((150, 23, 893, 25)\). After crossover, offspring \(X_{\text{off}}\) is \((150, 19, 893, 25)\) and offspring \(Y_{\text{off}}\) is \((145, 23, 825, 22)\). The mutation process
Hence, the offspring then randomly selects one component to mutate. In this process; if a random number within 0 and 1 is smaller than Pe is the probability of performing the mutation process; if a random number is smaller than Pe, then the process crossover proceeds. Similar to Pe, Pm is the probability of performing the mutation process; if a random number is smaller than Pm, then the process crossover proceeds. The higher Pm is, the faster the process performs a U-turn on the 3’ end. Besides, the length of each primer is <18 bp. SGD finds a solution that has the longest product size.

However, the solution does not satisfy the pair-complementary sequence and self-complementary constraint. None of the three design software can satisfy the specificity.

The product size of the solution obtained from the proposed algorithm is 303 bp. The melting temperatures are within 40–60°C. In addition, the primer pair found by the proposed algorithm also satisfies the complementary sequence. The temperature difference between the pair of primers is the highest among all the design software, but still within an acceptable range. It is time consuming to find the primer pair.

EXPERIMENTS

Materials

cDNA templates Two cDNA templates were used as our targets for PCR amplification, and subjected to primer designs in our experiments: Pseudomonas mendocina PHA synthase 1 (phaC1), PHA depolymerase (phaZ) and PHA synthase 2 (phaC2) genes, complete cds and Homo sapiens CDK2-associated protein 1 (CKD2AP1) coding DNAs (CDs) (GenBank acc. no. NM_004642; 523..870). One of the major differences between these two DNA sequences is their length.

Preparation of cDNA templates Total RNA was first extracted from the HeLa cell line. Reverse transcription into cDNAs were followed. The cDNAs were then subjected to being quality checked, quantified and adjusted to an appropriate concentration for further PCR reaction.

Dry dock experiments

The proposed algorithm was run on a Pentium 4, 1.5G Hz, 128 MB, Windows 2000 and jdk1.4.0 platform.

Results of the dry dock experiments The proposed algorithm was compared with three famous free primer design software (Primer3, GeneFisher and SGD). In Table 1, the primer pair, which the proposed algorithm finds, contains one similar restriction site on the forward primer. This similar restriction site can be corrected as ‘GGATCC’, so that the enzyme BamHI can be used and ‘CCTAGG’ (AvrII) on the reverse primer. Primer 3 finds a primer pair that satisfies the complementary sequence, but has the highest melting temperature. GeneFisher finds a primer pair whose product size is the longest, but it does not have specificity. SGD finds a primer pair, that contains the enzyme ‘AAATTT’ (DraI) on the forward primer and ‘GGGCC’ (ApaI) on the reverse primer, but it does not have specificity, and the complementary sequence is not satisfied.

The results of the second experiment are listed in Table 2. Primer 3 finds a primer pair that has a minimum temperature difference between the pair of primers. Besides, the primer pair also satisfies the complementary sequence. The melting temperatures of the two primers are around the optimal experimental temperature of 40–60°C. The PCR product amplified from primers designed by Primer 3 is 267 bp in length. GeneFisher finds a solution whose length is 252 bp. The melting temperatures are within 40–60°C. The primer pair satisfies the pair complementarity constraint, but it does not satisfy the self-complementary constraint, since the forward primer forms a U-turn on the 3’ end. Besides, the length of each primer is <18 bp. SGD finds a solution that has the longest product size.

However, the solution does not satisfy the pair-complementary sequence and self-complementary constraint. None of the three design software can satisfy the specificity.
Table 1. Comparisons of primer design among this study, primer3, GeneFisher and SGD for the first experiment

<table>
<thead>
<tr>
<th></th>
<th>The proposed algorithm</th>
<th>Primer 3</th>
<th>GeneFisher</th>
<th>SGD</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Forward (F) primer</strong></td>
<td>TTTTCACTGCTGATCTTCAAC</td>
<td>CTGACTAGAAACGTCGTA</td>
<td>CTCGAACTGAGAAGTCA</td>
<td>ACGGATTACGTAGCAGGA</td>
</tr>
<tr>
<td><strong>Reverse (R) primer</strong></td>
<td>ATCCGTCTAAGACTTTCACT</td>
<td>TCCGATTCTCAGCTCTTG</td>
<td>AGCCGATTACGTAGCAGGA</td>
<td>TTTGCCATTATCAGAGG</td>
</tr>
<tr>
<td><strong>Position (F)</strong></td>
<td>1814–1832</td>
<td>1968–1987</td>
<td>231–249</td>
<td>1890–1907</td>
</tr>
<tr>
<td><strong>Position (R)</strong></td>
<td>2905–2924</td>
<td>3064–3083</td>
<td>1413–1430</td>
<td>2987–3005</td>
</tr>
<tr>
<td><strong>Product size (bp)</strong></td>
<td>1112</td>
<td>1116</td>
<td>1200</td>
<td>1116</td>
</tr>
<tr>
<td><strong>Primer length (F/R)</strong></td>
<td>19/20</td>
<td>20/20</td>
<td>19/18</td>
<td>18/18</td>
</tr>
<tr>
<td><strong>Melting temperature</strong></td>
<td>54/56</td>
<td>60/62</td>
<td>56/54</td>
<td>51/51</td>
</tr>
<tr>
<td><strong>GC content (F/R)</strong></td>
<td>42/50</td>
<td>47/44</td>
<td>44</td>
<td></td>
</tr>
<tr>
<td><strong>Temperature difference (°C)</strong></td>
<td>2</td>
<td>2</td>
<td>0</td>
<td></td>
</tr>
<tr>
<td><strong>Specificity</strong></td>
<td>Yes</td>
<td>No</td>
<td>Yes</td>
<td>No</td>
</tr>
<tr>
<td><strong>Restriction site</strong></td>
<td>‘G’GTTACC’ (BamHI)/</td>
<td>N/A</td>
<td>N/A</td>
<td>‘AAATTT’ (DraI)</td>
</tr>
<tr>
<td></td>
<td>‘CCTAGG’ (AvrII)</td>
<td></td>
<td></td>
<td>‘GGGCC’ (Apel)</td>
</tr>
</tbody>
</table>

Table 2. Comparisons of primer design among this study, primer3, GeneFisher and SGD for the second experiment

<table>
<thead>
<tr>
<th></th>
<th>The proposed algorithm</th>
<th>Primer 3</th>
<th>GeneFisher</th>
<th>SGD</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Product size (bp)</strong></td>
<td>303</td>
<td>267</td>
<td>252</td>
<td>348</td>
</tr>
<tr>
<td><strong>Primer length (mer)</strong></td>
<td>18/18</td>
<td>18/18</td>
<td>16/17</td>
<td>21/18</td>
</tr>
<tr>
<td><strong>GC content (F/R)</strong></td>
<td>39/50</td>
<td>61/50</td>
<td>53</td>
<td>47</td>
</tr>
<tr>
<td><strong>Melting temperature</strong></td>
<td>50/50</td>
<td>58/54</td>
<td>52</td>
<td>56</td>
</tr>
<tr>
<td><strong>Temperature difference in pair (°C)</strong></td>
<td>4</td>
<td>2</td>
<td>2</td>
<td>3</td>
</tr>
<tr>
<td><strong>Specificity</strong></td>
<td>Yes</td>
<td>No</td>
<td>No</td>
<td>No</td>
</tr>
<tr>
<td><strong>3′-Terminus (F)</strong></td>
<td>C</td>
<td>A</td>
<td>A</td>
<td>C</td>
</tr>
<tr>
<td><strong>3′-Terminus (R)</strong></td>
<td>T</td>
<td>T</td>
<td>A</td>
<td>T</td>
</tr>
<tr>
<td><strong>Self-complementarity</strong></td>
<td>No</td>
<td>Yes</td>
<td>Yes</td>
<td>Yes</td>
</tr>
<tr>
<td><strong>Pair-complementarity</strong></td>
<td>No</td>
<td>No</td>
<td>Yes</td>
<td>Yes</td>
</tr>
</tbody>
</table>

that satisfies specificity and at the same time search for the primer pair that has a restriction site that can be recruited in the primers. However, the difference of the execution time among the proposed algorithm and the other design software is insignificant.

**Wet experiment**

**PCR amplifications** One primer pair, designed using our proposed algorithm, and based on CDK2AP1CDs (forward: 5′-ATGTCTTACAAACCGAAG-3′; reverse: 5′-CAGTCCTCTAGCGTGATT-3′) was used to perform the PCR reaction. The forward and reverse primers spanned from bases 1–18 and 285–303, respectively. A total of 15 µl reagent consisting of 50 ng HeLa cell cDNA, 10× reaction buffer 1.5 µl, dNTP (1.25 mM) 1.8 µl, forward and reverse primers (10 mM) 0.5 µl each, Taq polymerase (5 U/µl) 0.05 µl and deionized, distilled water, were mixed for PCR amplification. A hot start PCR program was set up as follows: 95°C for 5 min; 95°C for 45 s, 55°C for 45 s, 72°C for 45 s for 35 cycles;
The result of 1% agarose gel analysis in the second experiment (human CDK2AP1 complete cDNAs).

72°C for 10 min or final extension and then ramped to 4°C for 10 min.

Agarose gel electrophoresis and amplified CDK2AP1 PCR product

The amplified PCR products were analyzed in 1% agarose gel electrophoreses containing 0.1% ethidium bromide. Electrophoreses were conducted in 0.5× TBE buffer. After electrophoreses, the agarose gels were removed from the gel boxes and visualized under the illumination of UV light. Analysis of amplified CDK2AP1 PCR product (size = 303 bps) is shown in the middle lane of Figure 5.

The wet experimental results

Figure 5 shows the analysis of the second experimental PCR product, human CDK2AP1 complete cDNAs. In Figure 5, M shows a 100 bp ladder marker. Positive control (G3PDH) and negative control (water) are also included. At the middle of the gel is the amplified PCR product. Its length, which is ∼303 bp, can be clearly observed. The positive control result is at the left. Since there is only one band in positive control, the experimental result is acceptable. The negative control shows no bands in this experiment, hence it is certain that the experiment is not contaminated.

DISCUSSIONS

Primer design for PCRs is an extremely important issue in any molecular biology laboratory. That is, the quality of the primers decides whether or not the experiment will be successful. Many software for primer design based on different algorithms have been proposed in the last decade. However, none of these provided a tool to recruit specific restriction sites in a primer pair for further application in molecular cloning. In this study, we propose a method using GA that considers all these constraints. When compared with other primer design software, our proposed algorithm is able to find a feasible solution that follows all required properties in primer design. Besides, our two experiments show that no matter what the length of the DNA sequence is (long or short), the proposed algorithm is able to find a good solution.

One of the most critical properties in primer design is the specificity. Other properties for primer design, such as the length of the primers, can be achieved by using the fixing method–insertion, deletion and replacement. One method is able to modify the primer pair by extending or decreasing the length of the primer. This method is referred to as the insertion and deletion method. Another process is called the replacement method, and it relies on repairing the result. One can mutate the specific nucleic acid code of the primer to another nucleic acid code. For example, the nucleic acid code ‘A’ mutates to ‘T’, hence, it cannot be reduced just in order to speed up the evaluation process. In order to satisfy the specificity, the proposed algorithm loosens some constraints, like the 3′ termination and the GC content.

If the proposed algorithm ignores the specificity and the restriction site in the second experiment, then the execution time can be reduced to 1 s. However, in our experiment, although the proposed algorithm consumes time on handling the specificity and the restriction site, the running time is about the same as the execution time of the competing software without handling the specificity and the restriction site.

For the optimal problem, GA usually converges all solutions into one optimal solution in a short time by using mutation and crossover. Therefore, the redundant part can be reduced by examining the matrix. The proposed algorithm uses the matrix to record the position of the primer \( P_1 \) and the value of \( \text{Uni}(P_1) \). The restriction site checking process is similar to the specificity process. It also records the position of the primer \( P_1 \) and the value of the matched pattern’s length \( |P_m| \) to the enzyme sequence.

CONCLUSIONS

We presented a new algorithm using GA for primer design. The proposed algorithm met the design constraints such as melting temperatures, length, base composition, 3′-termini, repeated and self-complementary sequences and complementary sequences between members of a primer pair, especially for recruiting specific restriction sites and specificity. The sequencing result showed that there was a product whose length was approximately the solution that we expected, thus verifying the proposed method really can clip out the target sequence.
REFERENCES


